Algorithmic Approaches for Biological Data, Lecture #16

Katherine St. John

City University of New York American Museum of Natural History

30 March 2016

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Outline



• Networks & Graphs

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Outline



- Networks & Graphs
- Standard Representations: Adjacency Lists and Adjacency Matrices

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- Standard Representations: Adjacency Lists and Adjacency Matrices

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Reframing Biology Questions

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Problem Solving with Algorithms and Data Structures

K. St. John (CUNY & AMNH)

Algorithms #16

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• Graphs (networks)



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• Graphs (networks) have vertices (nodes)



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 - $V = \{V0, V1, V2, V3, V4, V5\}$



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- In example: G = (V, E) where:

$$V = \{V0, V1, V2, V3, V4, V5\}$$

$$E = \{(V0, V1, 5), (V1, V2, 4), (V2, V3, 9), (V3, V4, 7), (V4, V0, 1), (V0, V5, 2), (V5, V4, 8), (V3, V5, 3), (V5, V2, 1)\}$$



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Since edges have a direction, called a directed graph.



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- A cycle is a path that starts and ends at the same vertex.
 Example: (V5, V2, V3, V5) is a cycle.
- A graph with no cycles is called an acyclic graph.
- A directed graph with no cycles is called a directed acyclic graph (DAG).



	VO	V1	V2	V3	V4	V5
VO		5				2
V1			4			
V2				9		
V3					7	3
V4	1					
V5			1		8	

Problem Solving with Algorithms and Data Structures



• In Python, can use a list of lists or a numpy array.



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• In Python, can use a list of lists or a numpy array.

•	import numpy as np
	adjMatrix = np.zeros(6,6)
	adjMatix[0,1] = 5
	adjMatix[0,5] = 2
	adjMatix[1,2] = 4
	adjMatix[2,3] = 9
	adjMatix[3,4] = 7
	adjMatix[3,5] = 3
	adjMatix[4,0] = 1
	adjMatix[5,2] = 1
	adjMatix[5,4] = 8



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- import numpy as np adjMatrix = np.zeros(6,6) adjMatix[0,1] = 5 adjMatix[0,5] = 2 adjMatix[1,2] = 4 adjMatix[2,3] = 9 adjMatix[3,4] = 7 adjMatix[3,5] = 3 adjMatix[4,0] = 1 adjMatix[5,2] = 1 adjMatix[5,4] = 8
- Need to keep track of the node names separately.



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Advantages:



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- Advantages:
 - Checking if an edge occurs is quick.
 - Can check connectivity by matrix multiplication (explained in lab).
- Disadvantages:
 - ► Always the same size (n × n) even if there are few edges.



Problem Solving with Algorithms and Data Structures

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```
• import numpy as np
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adjList["V0"] = [("V1",5), ("V5",2)]
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adjList["V2"] = [("V3",9)]
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- < A > < B > < B >



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```

• Can look up each list of adjacencies in the dictionary using the vertex label as the key.

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1 - "V4"

adj = { V0:1 }

adj = { V2:1, V4:8

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V5

numMertices = 6

Advantages:

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Advantages:

 More space efficient for sparsely connected graphs





- Advantages:
 - More space efficient for sparsely connected graphs
- Disadvantages:





- Advantages:
 - More space efficient for sparsely connected graphs
- Disadvantages:
 - Could be costly to find adjacencies if a vertex has many neighbors.

In pairs/triples, represent the following graphs in the computer:



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• Lab today: connectivity & storing trees

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- Using networkx in lab today (for displaying graphs).



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- Email lab reports to kstjohn@amnh.org



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